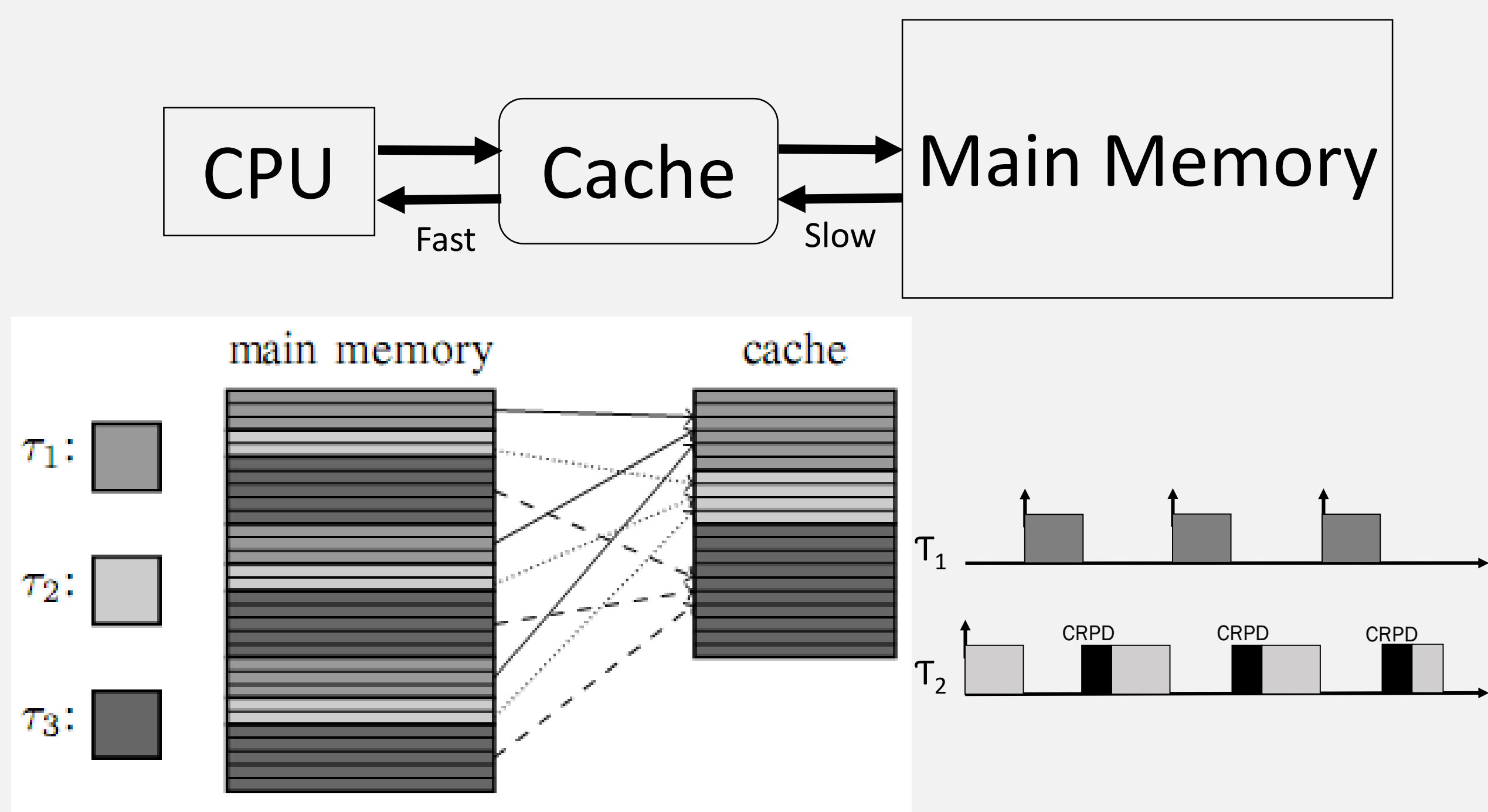


Cache Persistence Aware Response Time Analysis for Fixed Priority Preemptive Systems

1. Motivation

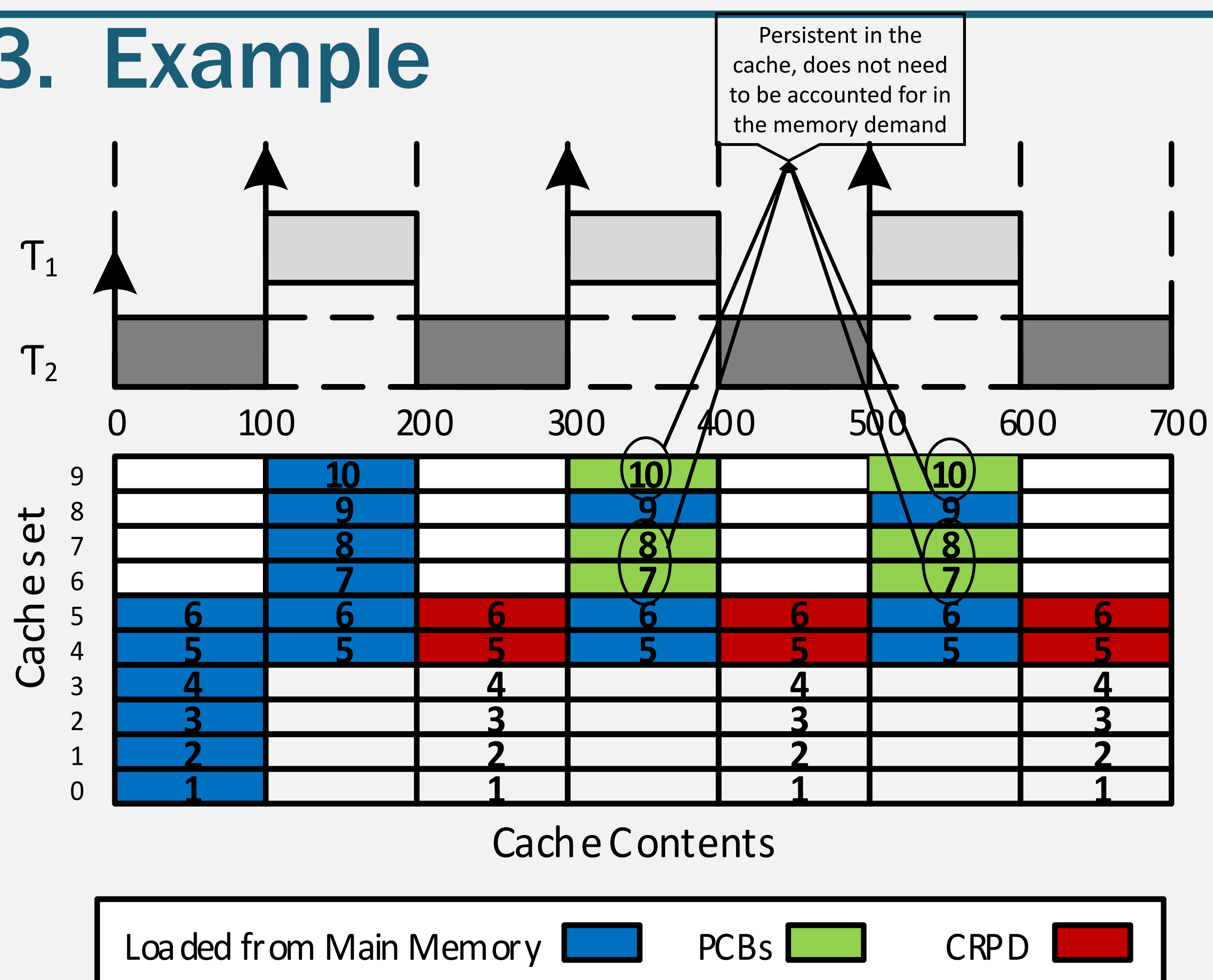
- Variations in WCET/WCRT due to cache hit/miss.
- Low priority tasks may need to account for **cache evictions** due to execution of high priority tasks (**CRPD**).
- State-of-the-art approaches for CRPD calculation **only** consider the **impact of high priority tasks** on memory demand of low priority ones.
- However, state-of-the-art **does not consider** the effect of low priority tasks on the memory demand of high priority tasks.



2. Contributions

- Preempting task can have content persisting in the cache between successive job executions.
- We introduce the concept of **cache persistence** in the context of WCRT analysis.
- We model the effect of **Persistent Cache Blocks (PCBs)** on the memory demand of preempting tasks.
- We account for the number of PCBs that can be evicted, i.e., **Cache Persistence Reload Overhead (CPRO)**

3. Example



References

[1] S. Altmeyer, R. Davis, C. Maiza et al., "Cache related pre-emption delay aware response time analysis for fixed priority pre-emptive systems," in RTSS'11. IEEE, 2011, pp. 261–271

4. Calculating CPRD and CPRO

4.1 ECB-Union and UCB-union Multi-set Approaches

- CRPD is usually calculated using **Evicting Cache blocks (ECBs)** of the preempting task and **Useful Cache Blocks (UCBs)** of the preempted task.
- ECB-union** approach considers ECBs of the preempting task T_j as well as all tasks in $hp(j)$.
- UCBs** of all tasks in $aff(i,j)$, i.e., $hp(i)$ and $lp(j)$ can be evicted by T_j .

$$CRPD_{i,j} = \max_{\forall k \in aff(i,j)} (|UCB_k \cap (\cup_{\forall j \in hp(j)} ECB_j)|)$$

- UCB-union multi-set** approach improve upon the ECB-union approach by additionally taking into account the **actual number of jobs** released by T_j and all tasks in $aff(i,j)$.

4.2 A Union approach for CPRO calculation

- Persistent Cache blocks (PCBs)** of a preempting task T_j executing during the response time of T_i can be evicted due to executions of tasks in $hp(i)$.
- Similar** formulation to ECB-union approach but considering PCBs of the **preempting task** T_j Instead of UCBs.
- ECBs of all tasks in $hp(i)$ can evict PCBs of T_j .

$$CPRO_{j,i} = PCB_j \cap (\cup_{\forall k \in hp(i)} ECB_k)$$

5. Improved WCRT Analysis

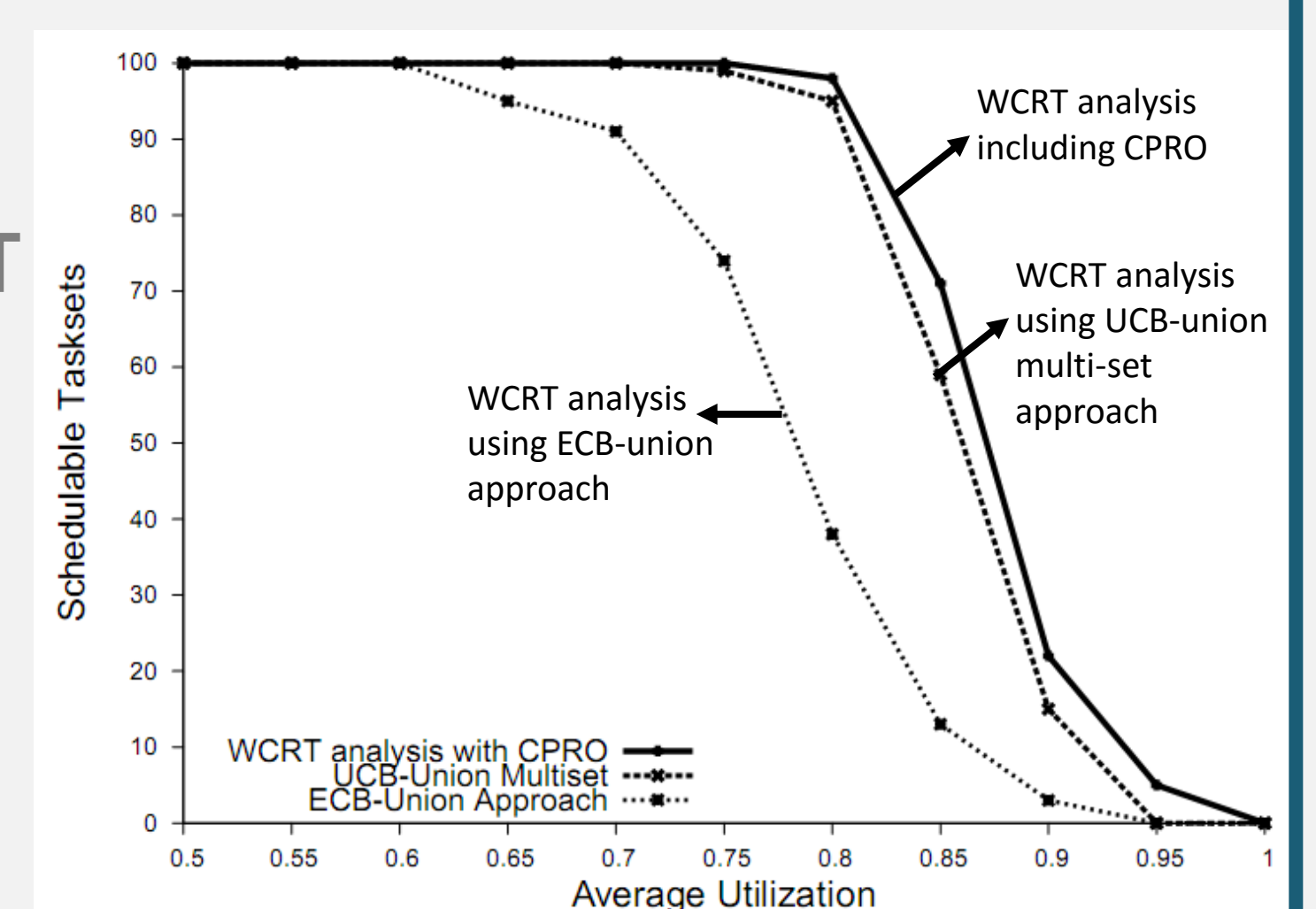
- Existing WCRT analysis for FPPS in the state-of-the-art only account for CRPDs.
- Our **Proposed WCRT analysis** incorporates for both CRPDs and CPRO, resulting in less pessimistic WCRT bounds.

$$R_i(t) = P_i + MD_i + \sum_{\forall j \in hp(i)} P_j + MD_j + \sum_{\forall j \in hp(i)} T_j \left[\frac{R_i}{T_j} - 1 \right] * (P_j + MD_j + CPRO_{j,i})$$

One job of T_j will load all its ECBs (max memory demand)
Accounting for evictions of UCBs of T_i
Accounting for evictions of PCBs of T_j
Memory Demand of T_j when all its PCBs are in the cache

5.1 Preliminary Results

- Initial results shows that **WCRT analysis with CPRO dominates** the ECB-union and UCB-union multi-set approaches.



6. Future Work

- In future, we plan to extend the analysis to set associative and data caches.
- Provide a less pessimistic multi-set approach to calculate CPRO.
- A combined approach to calculate both CRPD and CPRO.
- Extensive experimental evaluation using available benchmarks by varying different system parameters.